Orthomodular Algebraic Lattices and Combinatorial Posets

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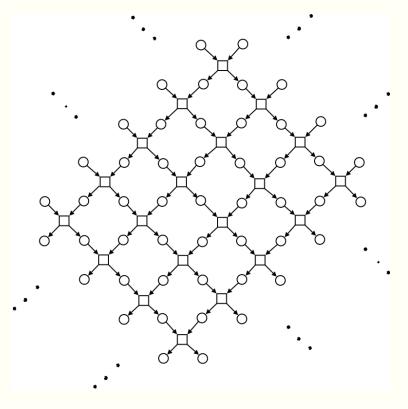
Partial orders and concurrent processes

 $P = (X, \leq)$

$$li = \leq \cup \geq \qquad co = (X \times X) \setminus li$$

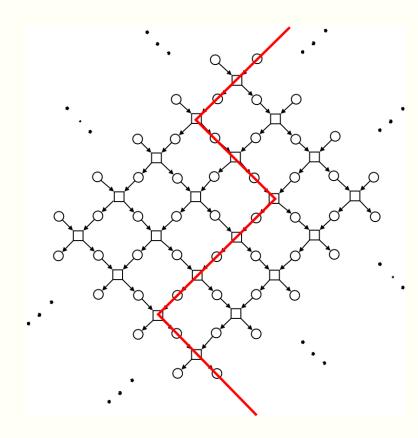
li and *co* are similarities (symmetric, non-transitive relations)

Combinatorial: $\leq = (\lessdot)^*$

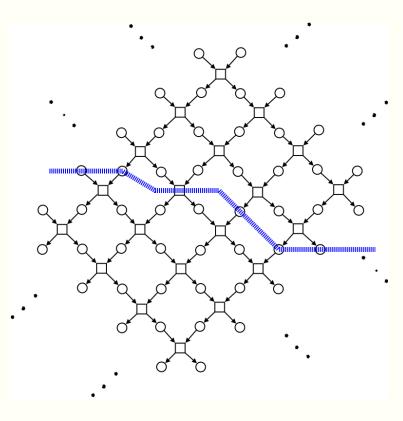


Lines and cuts

line: maximal clique of *li*



cut: maximal clique of $co \cup id_X$

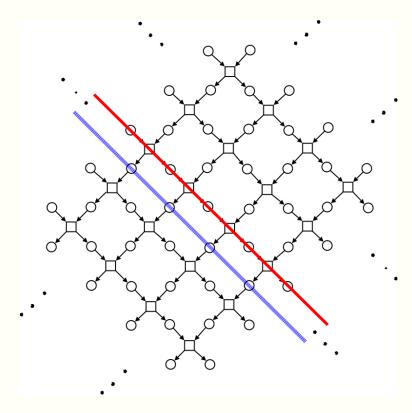


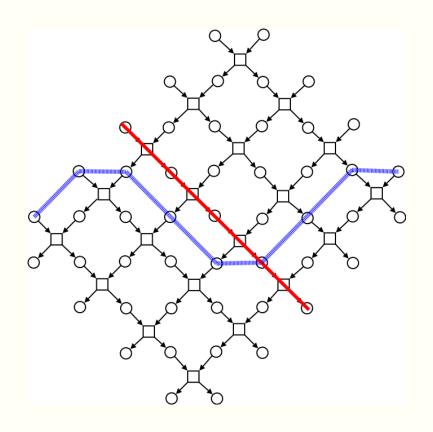
K-density and N-density

K-density: every line crosses every cut

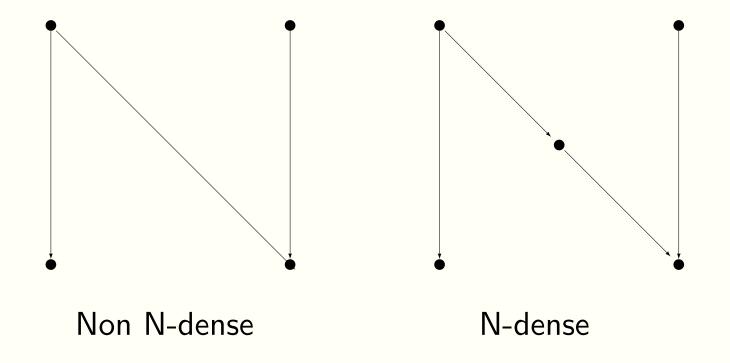
P is not K-dense

P is K-dense





N-density: a local form of K-density



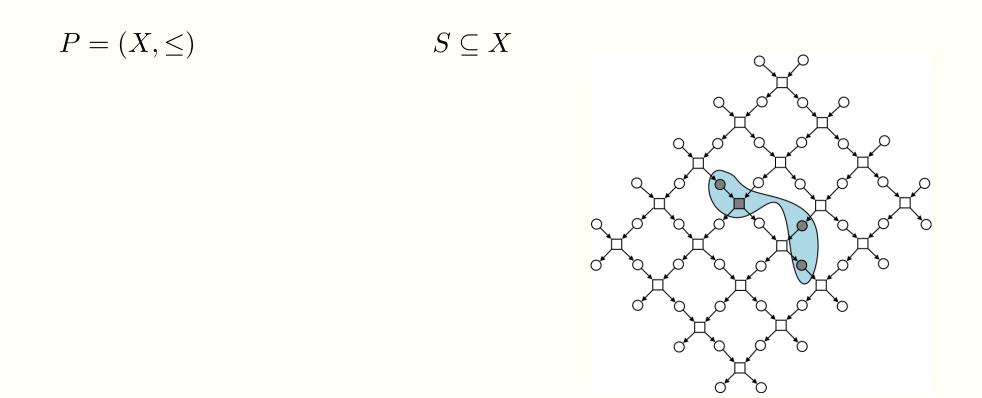
Closure operators

Closure operator on a set X

 $C: \mathbb{P}(X) \to \mathbb{P}(X)$

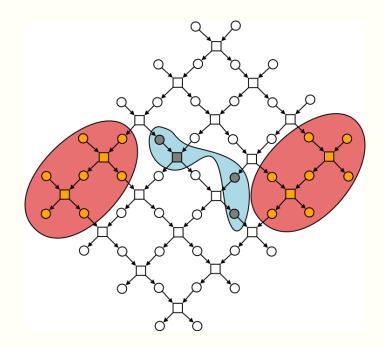
- $A \subseteq X$ $A \subseteq C(A)$
- $A, B \subseteq X$ $A \subseteq B \Rightarrow C(A) \subseteq C(B)$
- $A \subseteq X$ C(C(A)) = C(A)

A closure operator based on concurrency



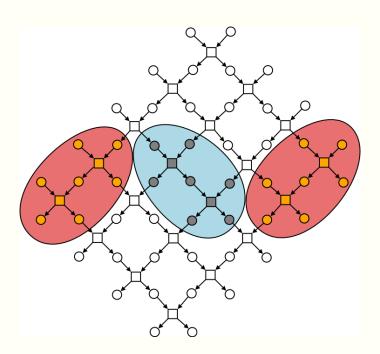
A closure operator based on concurrency

$$P = (X, \leq) \qquad \qquad S \subseteq X$$
$$S^{\perp} = \{x \in X \mid \forall y \in S : (x, y) \in co\}$$

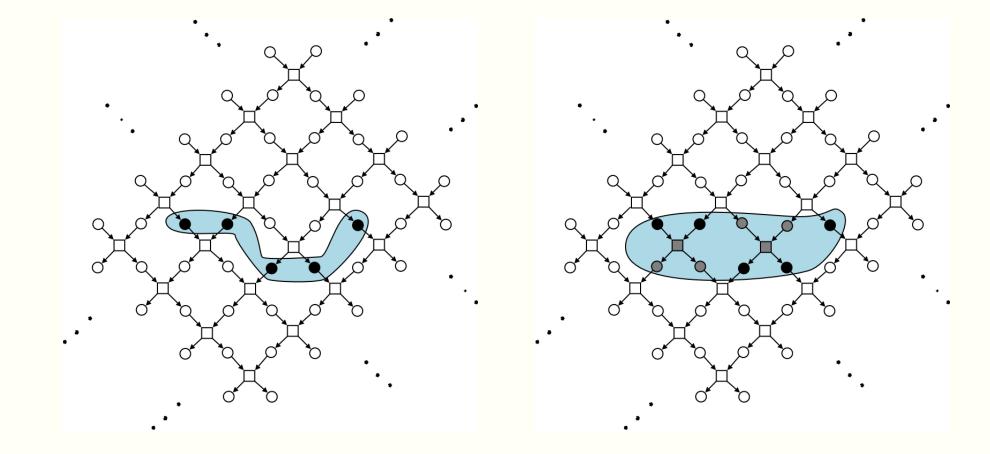


A closure operator based on concurrency

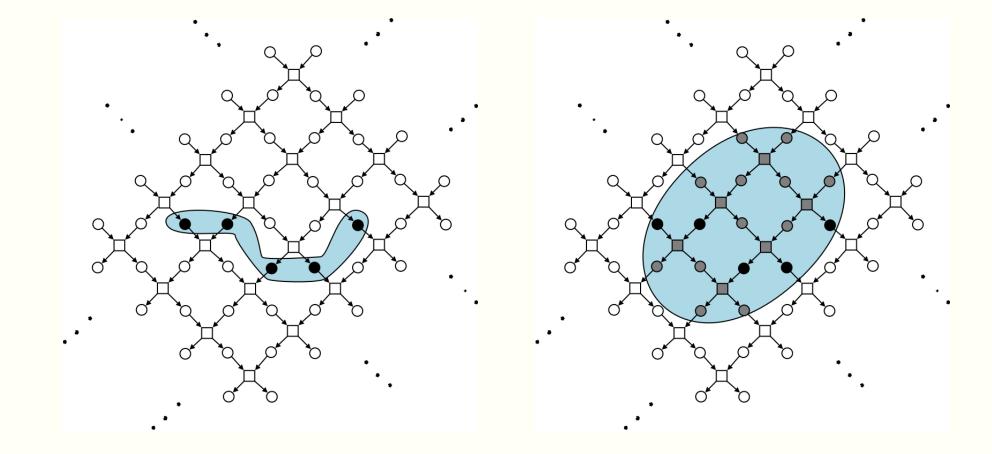
$$\begin{split} P &= (X, \leq) & S \subseteq X \\ S^{\perp} &= \{ x \in X \mid \forall y \in S : (x, y) \in co \} \\ (.)^{\perp \perp} : \mathbb{P}(X) \to \mathbb{P}(X) \\ S &= S^{\perp \perp} \Leftrightarrow S \text{ is closed} \end{split}$$



Causally closed subprocesses



Causally closed subprocesses



Lattices generated by concurrency

L(P) collection of closed sets

 $(.)^{\perp}$ applied to elements of L(P) is an orthocomplement

 $\mathcal{L}(P) = (L(P), \subseteq, \emptyset, X, (.)^{\perp})$

is an orthocomplete lattice

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is an orthocomplete lattice

Theorem If $\mathcal{P} = (P, \leq)$ is combinatorial, then $\mathcal{L}(P)$ is orthomodular if and only if $\mathcal{P} = (P, \leq)$ is N-dense.

Orthomodular lattices

An orthocomplete lattice

 $(L, \leq, 0, 1, (.)')$

is orthomodular if

$$S_1 \leq S_2$$
 implies $S_2 = S_1 \vee (S_2 \wedge S_1')$

Any orthomodular lattice can be seen as a family of partially overlapping Boolean algebras.

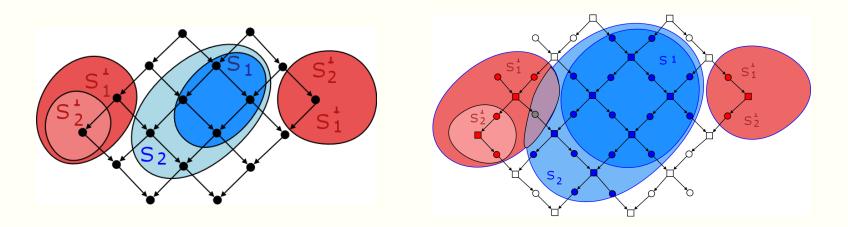
Lattices generated by concurrency

Theorem The closure operator $(.)^{\perp \perp}$ on a poset \mathcal{P} is algebraic if, and only if, \mathcal{P} is K-dense.

Lattices generated by concurrency on event structures

Elementary event structure (E, \sqsubseteq)

 $\mathcal{L}(E)$ is not orthomodular

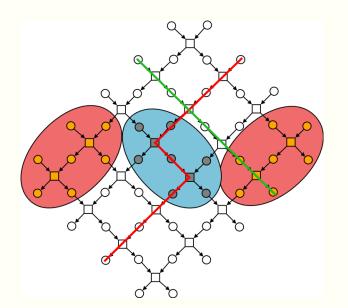


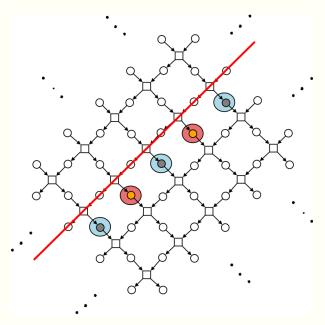
 $S_1 \subseteq S_2$ but $S_2 \neq S_1 \lor (S_2 \land S_1^{\perp})$

Properties of closed sets

If P is K-dense, then

S a closed set, λ a line: $\lambda\cap S\neq \emptyset$ or $\lambda\cap S^{\perp}\neq \emptyset$





Lines and two-valued states

 $\langle L, \leq, 0, 1, (.)' \rangle$ orthomodular lattice

Two valued state on \mathbf{L} $s: L \to \{0, 1\}$

•
$$s(1) = 1$$

•
$$s\left(\bigvee_{i} A_{i}\right) = \sum_{i} s(A_{i}) \quad A_{i} \in X, (i \neq j \Rightarrow A_{i} \leq A'_{j})$$

Lines and two-valued states

Let λ a line and $S \in L(P)$ a closed set

Define $\delta_{\lambda}(S) = 1$ if $S \cap \lambda \neq \emptyset$, and $\delta_{\lambda}(S) = 0$ otherwise

Theorem The map δ_{λ} is a two-valued state of L(P)

Conclusions

Ongoing work Closure operators defined on occurrence nets with conflicts (general unfoldings). Closure operators derived from the dependence relation, reflexive and irreflexive, closed sets as propositions.

Future developments Cyclic nets (cycloids).