

## THE CATEGORY OF AUTOMATA

The following definition, restricted to objects, is [Sipser, Definition 1.5].

**Definition 0.1.** An *automaton* consists of a 5-tuple  $\mathcal{A} = (Q, \Sigma, \delta, q_0, F)$  where  $Q$  and  $\Sigma$  are sets,  $q_0 \in Q$  is an element of  $Q$  and  $F \subset Q$  is a subset of  $Q$ , and where  $\delta: \Sigma \times Q \rightarrow Q$  is a function. The elements of  $Q$  are called *states* and the elements of  $\Sigma$  are called *letters*.

Suppose that  $\mathcal{A}' = (Q', \Sigma', \delta', q'_0, F')$  is another automaton. A *morphism of automata*, denoted  $(f, g): \mathcal{A} \rightarrow \mathcal{A}'$ , consists of a function  $f: Q \rightarrow Q'$  and a function  $g: \Sigma \rightarrow \Sigma'$  such that  $f(q_0) = q'_0$ ,  $f(F) \subset F'$  and the diagram

$$\begin{array}{ccc} Q \times \Sigma & \xrightarrow{\delta} & Q \\ f \times g \downarrow & & \downarrow f \\ Q' \times \Sigma' & \xrightarrow{\delta'} & Q' \end{array}$$

commutes. We denote the category of automata by **Automata**.

**Definition 0.2.** Let  $\mathcal{A} = (Q, \Sigma, \delta, q_0, F)$  be an automaton. We call the monoid  $\Sigma^*$  of sequences in  $\Sigma$  the *action monoid* for  $\mathcal{A}$ , and note that  $\delta$  gives  $Q$  the structure of a right  $\Sigma^*$ -set. For  $a \in \Sigma^*$  and  $q \in Q$ , we write  $q \cdot a$  to denote the result of acting  $a$  on  $q$ .

The *language* of  $\mathcal{A}$ , denoted  $L(\mathcal{A})$  is the subset of elements  $a \in \Sigma^*$  such that  $q_0 \cdot a \in F$ . Note that a morphism  $(f, g): \mathcal{A} \rightarrow \mathcal{A}'$  induces a morphism of sets  $L(\mathcal{A}) \rightarrow L(\mathcal{A}')$ , so that  $L: \mathbf{Automata} \rightarrow \mathbf{Sets}$  is a functor.

I'm not sure what the point of the following definition is, but it is natural to define.

**Definition 0.3.** Let  $\mathcal{A} = (Q, \Sigma, \delta, q_0, F)$  be an automaton. An element  $a \in \Sigma^*$  is said to *preserve acceptance* if, for all  $q \in F$  one has  $q \cdot a \in F$ . Note that the set of letters that preserve acceptance is a submonoid of  $\Sigma^*$ .

### REFERENCES

[Sipser] Sipser, M. *Introduction to the Theory of Computation* Second edition.